

The Fauser Monoid

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Diagrams

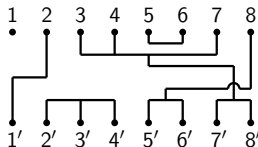
Let $k \in \mathbb{Z}^{>0}$, $K = \{1, \dots, k\}$ and $K' = \{1', \dots, k'\}$.

A **diagram** is a partition of $K \cup K'$.

Example

Take $k = 8$ and consider

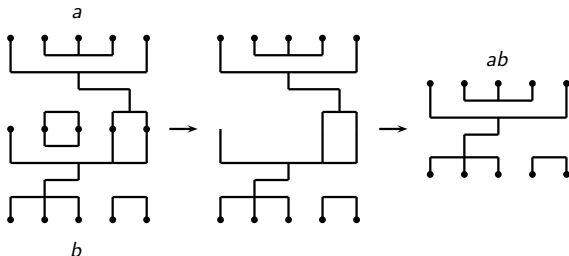
$\{\{1\}, \{2, 1'\}, \{3, 4, 7, 7', 8'\}, \{5, 6\}, \{8, 5', 6'\}, \{2', 3', 4'\}\}$.



- ▶ A diagram is **planar** if the edges can be drawn without crossing inside the rectangle bounding the vertices;
- ▶ **Transversal components** are edges that connect vertices in both rows;
- ▶ The **rank** of a diagram is the number of transversals it has.

Partition Monoid \mathcal{P}_k

Given two diagrams $a, b \in \mathcal{P}_k$, their product ab is formed pictorially as follows:



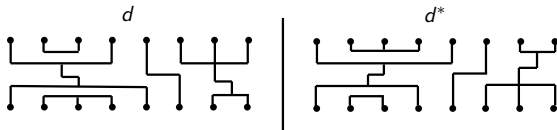
- ▶ place a on top of b ;
- ▶ remove the middle dots and stranded loops; and
- ▶ clip loose ends and collapse remaining loops.

The monoid of diagrams under this product is called the **partition monoid** \mathcal{P}_k .

\mathcal{P}_k is a regular $*$ -semigroup

Pictorially we obtain d^* by flipping d upside down.

Example



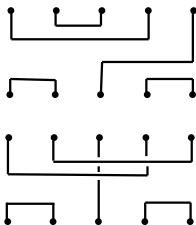
For each $d \in \mathcal{P}_k$:

- ▶ $d^{**} = d$ ($*$ is an involution);
- ▶ $(de)^* = e^* d^*$ ($*$ is an anti-homomorphism); and
- ▶ $dd^*d = d$ (regularity condition).

Jones Monoid \mathcal{J}_k and Brauer Monoid \mathcal{B}_k

- ▶ **Jones Monoid** \mathcal{J}_k consists of all planar matchings of $K \cup K'$;
- ▶ **Brauer Monoid** \mathcal{B}_k consists of all matchings of $K \cup K'$.

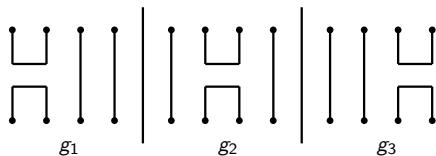
Examples



Diapsis Generators

Example

When $k = 4$ we have 3 **diapsis generators**:



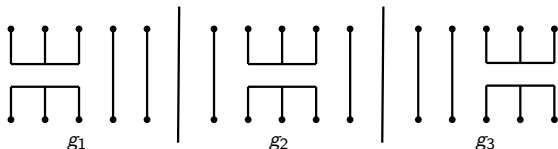
- ▶ \mathcal{J}_k is generated by diapsis generators (and $\text{id}_{\mathcal{P}_k}$);
- ▶ \mathcal{B}_k is generated by diapsis generators and S_k .

Triapsis Generators

Consider what happens when we replace the diapses in the generators of \mathcal{B}_k and \mathcal{J}_k with triapses.

Example

When $k = 5$ we have 3 **triapsis generators**:

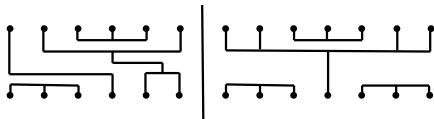


Obvious First Question: What diagrams are generated?

Triapsis Monoid \mathcal{T}_k and Fauser Monoid \mathcal{F}_k

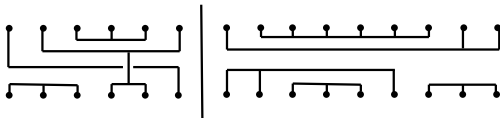
The **Triapsis Monoid** \mathcal{T}_k consists of $\underline{\text{id}_{\mathcal{P}_k}}$ and $d \in \mathcal{P}_k$ where:

- ▶ d is planar;
- ▶ there's at least one triapsis connecting consecutive points along the upper points, similarly along the lower points;
- ▶ for each edge $e \in d$, $|U(e)| \equiv |L(e)| \pmod{3}$.



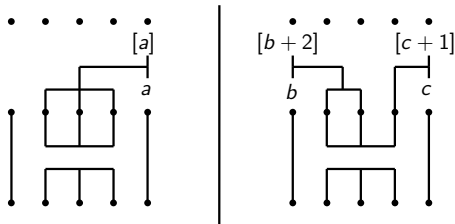
The **Fauser Monoid** \mathcal{F}_k consists of $\underline{S_k}$ and $d \in \mathcal{P}_k$ where:

- ▶ there's at least one triapsis along the upper points, similarly along the lower points;
- ▶ for each edge $e \in d$, $|U(e)| \equiv |L(e)| \pmod{3}$.



Characterising the elements of \mathcal{T}_k ($\langle g_1, \dots, g_{k-2}, \text{id}_{\mathcal{P}_k} \rangle = \mathcal{T}_k$)

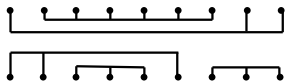
(\Rightarrow):) When showing $\langle g_1, \dots, g_{k-2}, \text{id}_{\mathcal{P}_k} \rangle \leq \mathcal{T}_k$, we show that $\overline{\mathcal{T}_k} \langle g_1, \dots, g_{k-2} \rangle \subseteq \mathcal{T}_k$.



(\Leftarrow):) To show that $\mathcal{T}_k \leq \langle g_1, \dots, g_{k-2}, \text{id}_{\mathcal{P}_k} \rangle$, we show how to decompose a diagram $d \in \mathcal{T}_k$ into a product of generators. We begin with $d = utl$, then decompose u , t and l .

Terminology Complications

- ▶ Triapsis Monoid is a better description of our generators than the elements of \mathcal{T}_k ;
- ▶ Can't call \mathcal{T}_k the planar version of \mathcal{F}_k ; and



- ▶ I'm not overly fond of referring to \mathcal{F}_k as the symmetric version of \mathcal{T}_k , plus if we can't think of descriptive names then we want to call \mathcal{F}_k the Fauser monoid.

Why Fauser?

Cardinality of \mathcal{F}_k

Let

- ▶ $N(n, t_1, \dots, t_k)$ be the number of ways to place t_1 triapses, ..., t_k $3k$ -apses along n points;
- ▶ $N(n, t) = \sum_{3t_1 + \dots + 3kt_k = t: t_i > 0} N(n, t_1, \dots, t_k)$ be the number of ways to use t of n dots with non-transversals; and
- ▶ $T(n_1, n_2)$ be the number of ways to feasibly connect n_1 points to n_2 points with just (feasible) transversals.

We have the following recurrence relations for N and T :

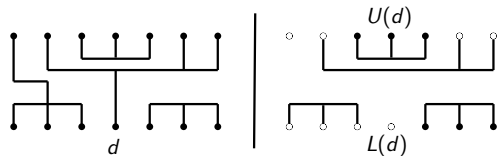
- ▶ $N(n, 0, \dots, 0) = 1$;
- ▶ $N(n, t_1, \dots, t_k) = N(n-1, t_1, \dots, t_k) + \sum_{i: t_i > 0} \binom{n-1}{3i-1} N(n-3i, t_1, \dots, t_i-1, \dots, t_k)$;
- ▶ $T(0, 0) = 1$, $T(n_1, n_2) = T(n_2, n_1)$, $T(n_1, 0) = 0$, $n_1 > 0$;
- ▶ $T(n_1, n_2) = \sum_{(n'_1, n'_2) \leq (n_1, n_2): n'_1 \equiv n'_2 \pmod{3}} \binom{n_1-1}{n'_1-1} \binom{n_2}{n'_2} T(n_1 - n'_1, n_2 - n'_2)$.

$$|\mathcal{F}_k| = n! + \sum_{u=1}^{\lfloor n/3 \rfloor} \sum_{l=1}^{\lfloor n/3 \rfloor} N(n, 3u) N(n, 3l) T(n-3u, n-3l).$$

Patterns

A **pattern** p is a partition of K (or K') with a two-tone vertex colouring, where the colour of a vertex indicates whether the edge connected to it is a transversal or non-transversal component.

Example



Hence we can break a diagram $d \in \mathcal{P}_k$ into its **upper pattern** $U(d)$ and **lower pattern** $L(d)$.

\mathcal{T}_k -admissibility and \mathcal{T}_k -compatibility

- ▶ a pattern p is \mathcal{T}_k -**admissible** if $\exists d \in \mathcal{T}_k$ with $U(d) = p$; and
- ▶ \mathcal{T}_k -admissible p, q are \mathcal{T}_k -**compatible** if $\exists d \in \mathcal{T}_k$ with $U(d) = p$ and $L(d) = q$.
(d is unique, which we denote by $\delta(p, q)$)

Properties

- ▶ \mathcal{T}_k -compatibility is an equivalence relation; and
- ▶ $\delta(p, q)\delta(q, r) = \delta(p, r)$.

Characterising \mathcal{T}_k -admissibility and \mathcal{T}_k -compatibility

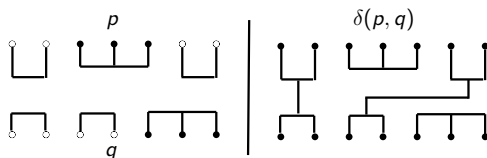
A pattern p is \mathcal{T}_k -admissible iff:

- ▶ it is planar;
- ▶ it has at least one triapsis; and
- ▶ each non-transversal component has cardinality divisible by 3.

\mathcal{T}_k -admissible patterns p, q are \mathcal{T}_k -compatible iff:

- ▶ $rank(p) = rank(q)$; and
- ▶ the cardinalities of *matched* transversal components are congruent mod 3.

Example



Green's Relations

Definition

For $a, b \in S$:

- ▶ $\mathcal{R} = \{(a, b) \in S^2 : aS^1 = bS^1\}$;
- ▶ $\mathcal{L} = \{(a, b) \in S^2 : S^1a = S^1b\}$;
- ▶ $\mathcal{H} = \mathcal{L} \cap \mathcal{R}$; and
- ▶ $\mathcal{J} = \{(a, b) \in S^2 : S^1aS^1 = S^1bS^1\}$;

Theorem (Howie)

If $T \leq S$ is regular then Green's \mathcal{L} , \mathcal{R} and \mathcal{H} relations are just their respective restrictions on T . I.e. $\mathcal{L}^T = \mathcal{L}^S \cap T^2$.

Green's Relations on \mathcal{P}_k

Theorem (Wilcox)

For $a, b \in \mathcal{P}_k$:

- ▶ $a\mathcal{R}b$ iff $U(a) = U(b)$;
- ▶ $a\mathcal{L}b$ iff $L(a) = L(b)$;
- ▶ $a\mathcal{H}b$ iff $U(a) = U(b)$ and $L(a) = L(b)$; and
- ▶ $a\mathcal{J}b$ iff $\text{rank}(a) = \text{rank}(b)$.

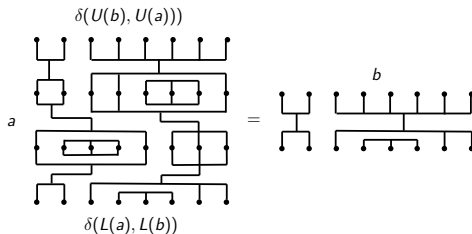
Green's \mathcal{J} Relation on \mathcal{T}_k

Theorem

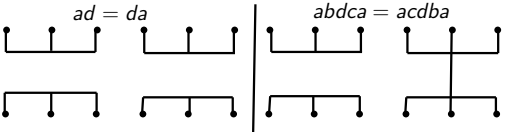
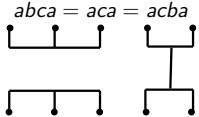
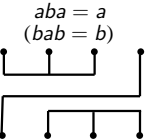
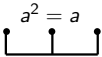
For $a, b \in \mathcal{T}_k$, $a \mathcal{J} b$ iff ' $U(a)$ and $U(b)$ are \mathcal{T}_k -compatible'.

(\Rightarrow) : ' $\text{rank}(a) = \text{rank}(ab)$ ' \Rightarrow ' $U(a)$ and $U(ab)$ are \mathcal{T}_k -compatible'.

$$\begin{aligned}
 \underline{(\Leftarrow)}: [\delta(U(b), U(a)) \cdot a] \cdot \delta(L(a), L(b)) &= \delta(U(b), L(a)) \cdot \delta(L(a), L(b)) \\
 &= \delta(U(b), L(b)) \\
 &= b
 \end{aligned}$$



Presentation of \mathcal{T}_k



Presentation of \mathcal{T}_k

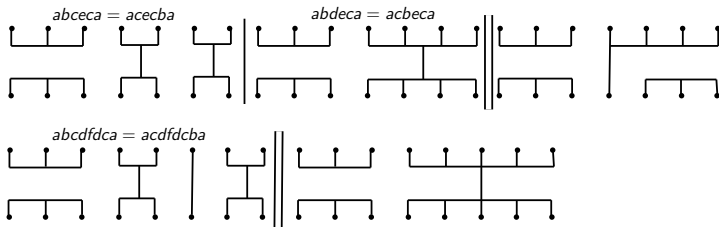
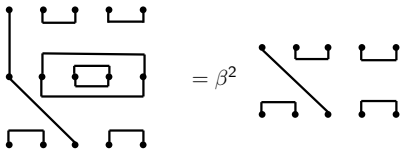


Diagram Algebras

- ▶ For $\beta \in \mathbb{C}$, the **partition algebra** $\mathbb{C}^\beta[\mathcal{P}_k]$ is the semigroup algebra $\mathbb{C}[\mathcal{P}_k]$ with multiplication $d * d' = \beta^r(dd')$ where r is the number of blocks removed when forming dd' .



- ▶ $\mathbb{C}^\beta[\mathcal{B}_k]$ and $\mathbb{C}^\beta[\mathcal{F}_k]$ are defined analogously.

Schur-Weyl Duality

Schur-Weyl duality tells us (amongst other things) that:

- ▶ $GL_n(\mathbb{C})$ and $\mathbb{C}[S_k]$ have commuting actions on $V^{\otimes k}$;
- ▶ Each action generates the full centraliser of the other
ie. $\text{End}_{GL_n(\mathbb{C})}(V^{\otimes k}) = \mathbb{C}[S_k]$ and $\text{End}_{\mathbb{C}[S_k]}(V^{\otimes k}) = GL_n(\mathbb{C})$.

A number of analogous dualities are known, for example between:

- ▶ $O_n(\mathbb{C}) \subseteq GL_n(\mathbb{C})$ and $\mathbb{C}^\beta[\mathcal{B}_k] \supseteq \mathbb{C}[S_k]$ (Brauer);
- ▶ $S_n \subseteq O_n$ and $\mathbb{C}^\beta[\mathcal{P}_k] \supseteq \mathbb{C}^\beta[\mathcal{B}_k]$ (Martin);
- ▶ IS_n and $\mathbb{C}[I_k^*]$ (Kudryavtseva & Mazorchuk).

We are hoping to find a subgroup of $GL_n(\mathbb{C})$ which is in a Schur-Weyl type duality with $\mathbb{C}^\beta[\mathcal{F}_k]$.

Suggestions and/or Questions?